

# Modeling Systems and Symbolic Encodings

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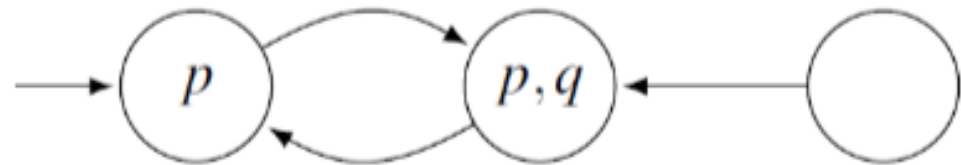
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# Reactive System

- **Set of States**
- Initial State(s)
- State Update Rules
  - aka *Transition Function* (deterministic case)
  - aka *Transition Relation* (non-Deterministic)

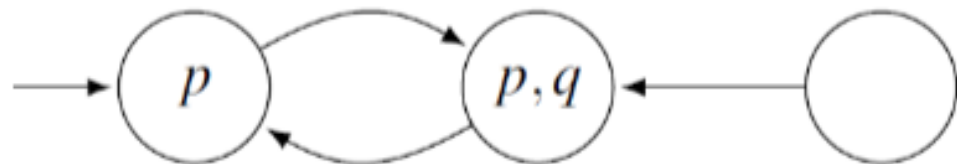
# Kripke Structure

- Formalization of Digital System
  - Finite Set of States  $S$
  - Set of Initial States  $I \subseteq S$
  - Transition Relation  $T \subseteq S \times S$
  - Labeling Function  $L: S \rightarrow 2^{AP}$ 
    - “AP”: Set of Atomic Propositions



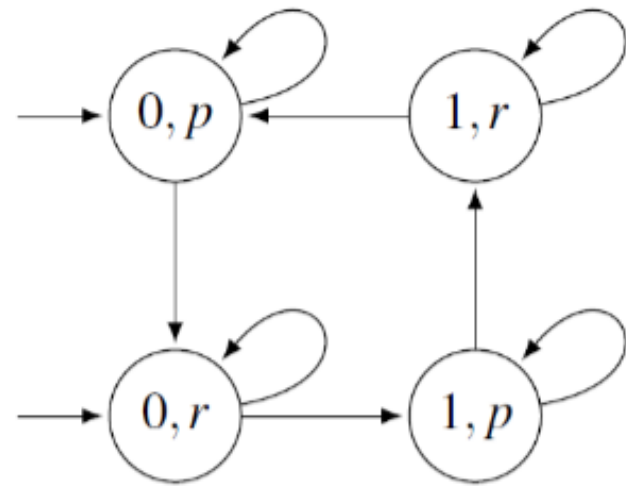
# Kripke Structure

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  - Transition Relation  $T \subseteq S \times S$
  - Labeling Function  $L: S \rightarrow 2^{AP}$ 
    - “AP”: Set of Atomic Propositions
    - $AP = \{p, q\}$
    - $L = \{s_1 \rightarrow \{p\}, s_2 \rightarrow \{p, q\}, s_3 \rightarrow \emptyset\}$



# Kripke Structure

- **Light Switch Example:**
  - Initially, the light is off. Once a button is pressed, the light is turned on. To turn the light off, the button has to be released and pressed again.
- **Model the light switch as Kripke structure**
  - **Labels:**
    - 1 ... the light is on
    - 0 ... the light is off.
    - r ... the button is released
    - p ... button is pressed



# Symbolic Encoding

- Systems have huge state spaces / number of transitions
- Therefore,
  - Symbolically encode sets
  - Perform set operations symbolically

# Boolean Function

- $f: \mathbb{B}^n \mapsto \mathbb{B}$ 
  - $\mathbb{B} = \{True, False\}$
- Possible Representations
  - Propositional Formula
  - Truth Table
  - Binary Decision Diagram

# Finite Sets vs. Boolean Functions

## ■ Set:

- Finite Domain of Objects

- Element

- “Non-Element”

## ■ Boolean Function:

- Finite Domain of Truth Assignments

- Map to True

- Map to False





# Characteristic Function

- Function  $f$  represents set  $A$ 
  - $\llbracket f \rrbracket = A$
- $a \in A \iff f(Enc(a)) = \top$ 
  - Elements give True
  - Non-Elements give False

# Encoding Elements

- Domain: e.g.  $D = \{Austria, Germany, Spain, Italy\}$ 
  - $\#Vars = \lceil \lg(|D|) \rceil$

Element	Encoding	
	$x_1$	$x_0$
Austria	0	0
Germany	0	1
Spain	1	0
Italy	1	1

# Example

- $f = \neg x_0 \wedge \neg x_1$

$$\llbracket f \rrbracket = \{Austria\}$$

- $g = \neg x_0$

$$\llbracket g \rrbracket = \{Austria, Spain\}$$

Element	Encoding	
	$x_1$	$x_0$
Austria	0	0
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Spain	1	0
Italy	1	1

# Example

Element	Encoding	
	$x_1$	$x_0$
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Element	Encoding	
	$x_1$	$x_0$
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- Which encoding gives the shorter characteristic function for the set  $B = \{\text{Germany, Spain}\}$ ?

$$f = x_1$$

$$g = x_1 \oplus x_0$$

# Example

1. Consider the set  $A = \{Apple, Banana, Orange, Pear\}$ . Find a symbolic encoding for its elements. Using this encoding, determine the characteristic functions of the sets  $B = \{Apple, Pear\}$ ,  $C = \{Orange\}$ , and  $D = \{\}$ .

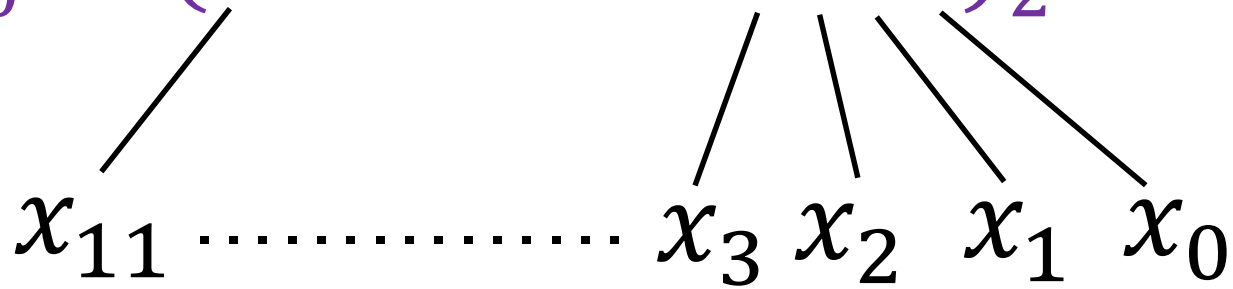
Element	Encoding	
	$x_1$	$x_0$
Apple	0	0
Banana	0	1
Orange	1	0
Pear	1	1

$$\begin{aligned}b &= (\neg x_1 \wedge \neg x_0) \vee (x_1 \wedge x_0) \\c &= x_1 \wedge \neg x_0 \\d &= \perp\end{aligned}$$

# Encoding Natural Numbers

- Binary Representation
- Domain  $D$ : Usually Power of 2
  - E.g.:  $D = \{x \in \mathbb{N} \mid x < 2^{12}\}$

$$(457)_{10} = (0001\ 1100\ 1001)_2$$



# Example

- Domain:  $A = \{x \in \mathbb{N} \mid 0 \leq x \leq 1023\}$ 
  - 10 bit binary representation  $x_9x_8 \dots x_0$
- $B = \{x \in A \mid x < 512\}$ ,  $b = ?$
- $C = \{x \in A \mid 256 \leq x < 768\}$ ,  $c = ?$

256	010...0	}	$c = (\neg x_9 \wedge x_8) \vee$ $(x_9 \wedge \neg x_8)$
511	011...1		
512	100...0	}	
767	101...1		

# Symbolic Operations

- **Intersection:**  $F \cap G =$   
 $= \llbracket f \wedge g \rrbracket$
- **Union:**  $F \cup G =$   
 $= \llbracket f \vee g \rrbracket$
- **Difference:**  $F \setminus G =$   
 $= \llbracket f \wedge \neg g \rrbracket$



# Symbolic Relations

- **Equality:**  $F = G \Leftrightarrow$

$$\Leftrightarrow f \equiv g$$

- **Subset:**  $F \subseteq G \Leftrightarrow$

$$\Leftrightarrow f \rightarrow g$$

# Example

- Domain:  $A = \{x \in \mathbb{N} \mid 0 \leq x \leq 1023\}$

- $b = \neg x_9$        $c = (\neg x_9 \wedge x_8) \vee (x_9 \wedge \neg x_8)$

- $D = B \cup C$

$$d = \neg x_9 \vee [(\neg x_9 \wedge x_8) \vee (x_9 \wedge \neg x_8)] = \neg x_9 \vee (x_9 \wedge \neg x_8)$$

- $E = B \cap C$

$$e = \neg x_9 \wedge [(\neg x_9 \wedge x_8) \vee (x_9 \wedge \neg x_8)] = \neg x_9 \wedge x_8$$

- $F = A \setminus E$

$$f = T \wedge \neg[\neg x_9 \wedge x_8] = x_9 \vee \neg x_8$$

$$|A| = 1024$$

$$|E| = 256$$

$$|F| = 768$$

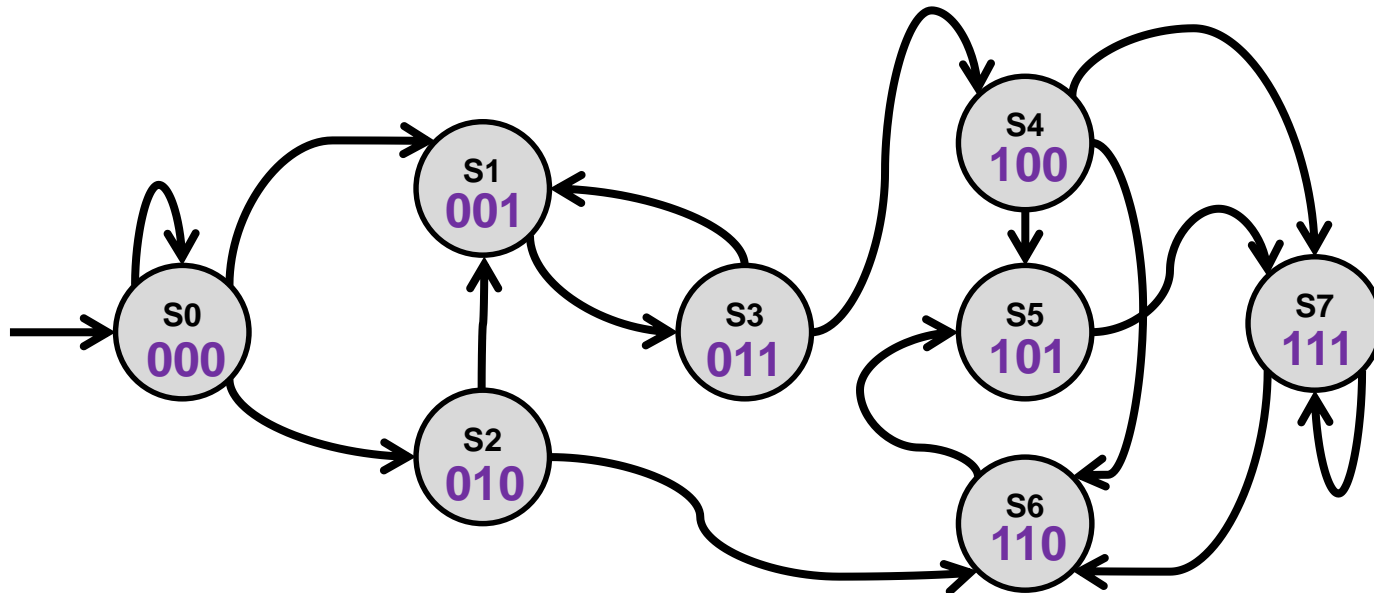
$$|B| = 512$$

# Symbolic Encoding of Kripke Structures

- Systems have huge state spaces / number of transitions
- Therefore:
  - Symbolically encode sets of states
  - Symbolically encode sets of transitions
  - Perform set operations symbolically

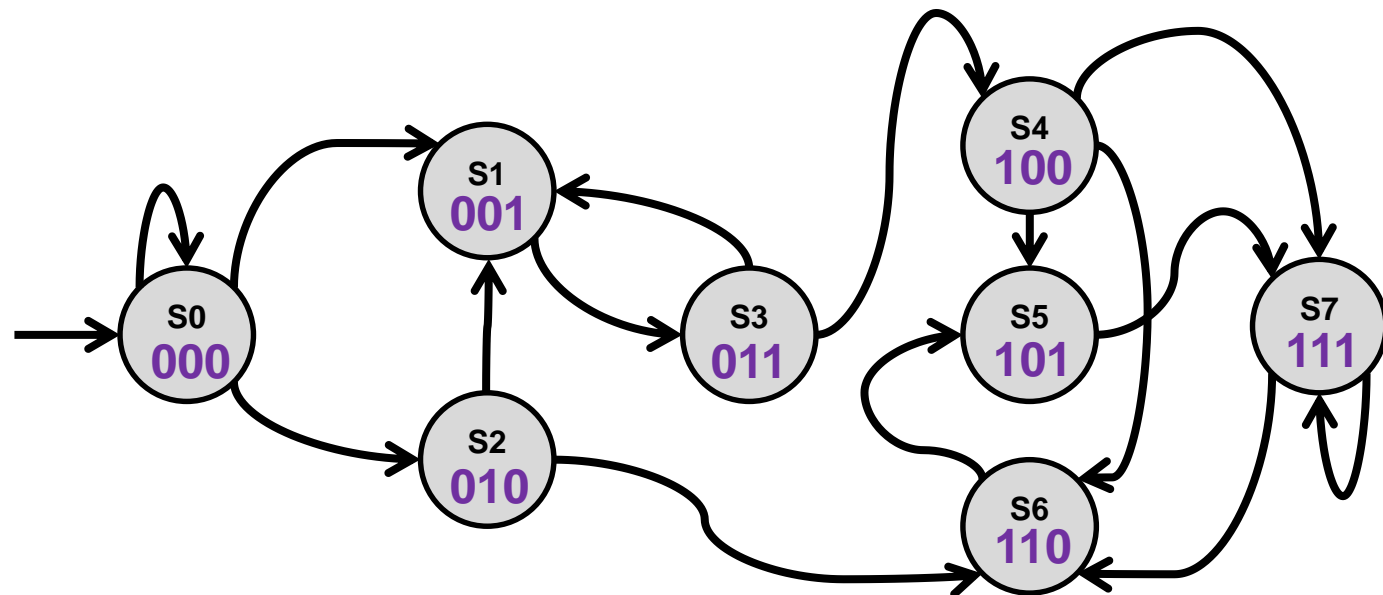
# Encoding Kripke Structures – States

- Use variables  $V = \{v_1, \dots, v_n\}$  for **binary representations** of states



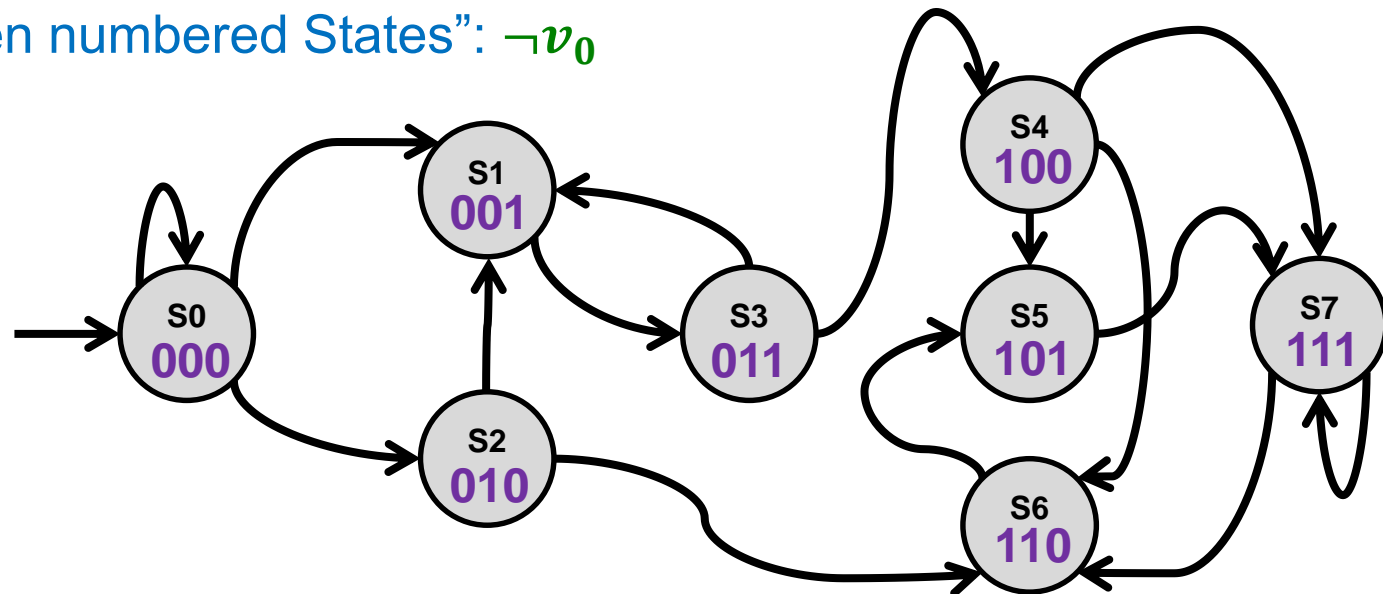
# Encoding Kripke Structures – States

- Single State
  - Labeling  $v_2v_1v_0$ 
    - E.g. State  $S_2$ :  $\neg v_2 \wedge v_1 \wedge \neg v_0$



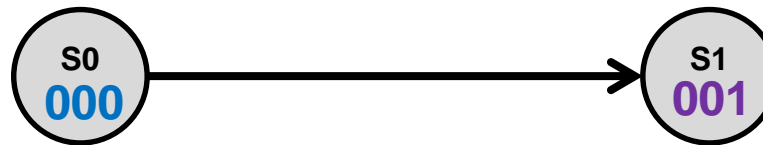
# Encoding Kripke Structures – States

- Single State
  - Labeling  $v_2v_1v_0$ 
    - E.g. State  $S_2 : \neg v_2 \wedge v_1 \wedge \neg v_0$
- Sets
  - E.g. “Even numbered States”:  $\neg v_0$



# Encoding Kripke Structures – Edges

- “Duplicate” Variables
  - “Present” State variables  $v_0, v_1, v_2, \dots$
  - “Next” State variables  $v'_0, v'_1, v'_2, \dots$

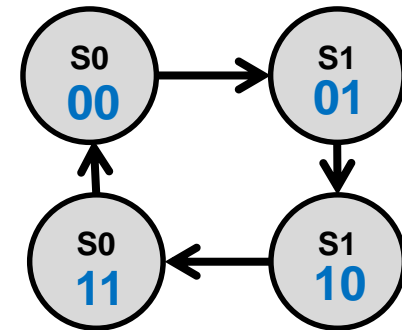


$$\neg v_2 \wedge \neg v_1 \wedge \neg v_0 \quad \wedge \quad \neg v'_2 \wedge \neg v'_1 \wedge v'_0$$

# Symbolic encoding for given transition-relation

- Draw the Kripke Structure for the following transition relation:

$$\begin{aligned} & (\neg v_1 \wedge \neg v_0 \wedge \neg v'_1 \wedge v'_0) \vee \\ & (\neg v_1 \wedge v_0 \wedge v'_1 \wedge \neg v'_0) \vee \\ & (v_1 \wedge \neg v_0 \wedge v'_1 \wedge v'_0) \vee \\ & (v_1 \wedge v_0 \wedge \neg v'_1 \wedge \neg v'_0) \end{aligned}$$





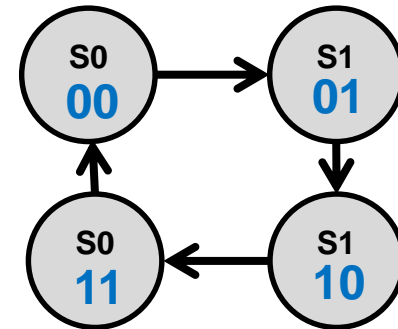
# Transition Function/Relation

- Union of all Edges
  - Disjunction
  - Good for sparse sets of edges
- $[[T]] \setminus \{missing\ edges\}$ 
  - Good for dense sets of edges
- Recognize Patterns
  - E.g. even numbered states have edges to (all) odd numbered states
  - $\neg x_0 \wedge x'_0$

# Example

- Symbolic encoding for given transition-relation?

$$\begin{aligned} & (\neg x_1 \wedge \neg x_0 \wedge \neg x'_1 \wedge x'_0) \vee \\ & (\neg x_1 \wedge x_0 \wedge x'_1 \wedge \neg x'_0) \vee \\ & (x_1 \wedge \neg x_0 \wedge x'_1 \wedge x'_0) \vee \\ & (x_1 \wedge x_0 \wedge \neg x'_1 \wedge \neg x'_0) \end{aligned}$$



# Summary

- Symbolic Sets
  - Encodings, Operations
- Kripke Structures
  - Graphs that represent digital systems, Symbolic Encoding